



DEPARTMENT OF MATHEMATICS COMPUTER SCIENCE & ENGINEERING
TECHNOLOGY

Introduction to Graph Polynomials

by

Shanee' Hamilton

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APPROVED BY

Julian D. Allagan, Ph.D.
Committee Chair

Weizheng Gao, Ph.D.
Committee Member

Gabriela H. Del Villar, Ph.D.
Committee Member

Committee member

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ABSTRACT

With graph polynomials being a fairly new but intricate realm of graph theory, I will begin with a brief historical background and progress to elucidate each polynomial's unique characteristics and mathematical underpinnings. Through illustrative examples, the paper elucidates the practical applications of these graph polynomials, showcasing their efficacy in real-world scenarios. My research contributes to the broader understanding of graph polynomials and inspires further research in the intersection of mathematics and technology.

DEDICATION

This thesis is dedicated to my parents Sang Sr. and Shirley Hamilton. Education has always been number one and I thank you all for instilling that in me and never giving up on me no matter what path I have chose along the way. To my children: ShaTreece, Donsha' Jr., and Elyse, thank you for being my inspiration. I would also like to dedicate this thesis to my twin, Janee'. She has always supported me on my journey. To my brother Deon, this one for you Deebo. I did it! Lastly, to close family and friends that have always believed in me no matter what I dedicate this to you too!

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graph	chromatic polynomial
barbell graph	$\frac{(z)_n^2 (z-1)}{z}$
book graph $S_{n+1} \square P_2$	$(z-1) z (z^2 - 3z + 3)^n$
centipede graph	$(z-1)^{2n-1} z$
complete graph K_n	$(z)_n$
cycle graph C_n	$(-1)^n (z-1) + (z-1)^n$
gear graph	$z [z - 2 + (3 - 3z + z^2)^n]$
helm graph	$z [(1-z)^n (z-2) + (z-2)^n (z-1)^n]$
ladder graph $P_2 \square P_n$	$(z-1) z (z^2 - 3z + 3)^{n-1}$
ladder rung graph $n P_2$	$z^n (z-1)^n$
Möbius ladder M_n	$-1 + (1-z)^n - (3-z)^n + (-1-z)^n + (3-z)^n z + (3 + (-3+z)z)^n$
pan graph	$(z-1)^{n+1} + (-1)^n (z-1)^2$
path graph P_n	$z (z-1)^{n-1}$
prism graph Y_n	$1 + [z(z-3) + 3]^n + z [(1-z)^n + (3-z)^n + z - 3] - (1-z)^n - (3-z)^n$
star graph S_n	$z (z-1)^{n-1}$
sun graph	$(z)_n (z-2)^n$
sunlet graph $C_n \odot K_1$	$(z-1)^{2n} - (1-z)^{n-1}$
triangular honeycomb rook graph	$\prod_{k=1}^n [(z)_k]^n$
web graph	$z [(1-z)^n + (3-z)^n + z - 3] (z-1)^n + (z-1)^n - [-(z-3)(z-1)]^n - [-(z-1)^2]^n + [(z-1)((z-3)z + 3)]^n$
wheel graph W_n	$z [(z-2)^{n-1} - (-1)^n (z-2)]$

The following table summarizes clique polynomials for some common classes of graphs.

graph	$C(x)$
Andrásfai graph A_n	$1 + \frac{1}{2}(3n-1)x(n x + 2)$
antiprism graph	$\begin{cases} (2x+1)(4x^2+6x+1) & \text{for } n=3 \\ nx^2+nx+1 & \text{otherwise} \end{cases}$
barbell graph	$-1+x^2+2(1+x)^n$
book graph $S_{n+1} \square P_2$	$(3n+1)x^2+x^2+2(n+1)x+1$
cocktail party graph $K_{n \times 2}$	$(1+2x)^n$
complete bipartite graph $K_{m,n}$	$(1+mx)(1+nx)$
complete graph K_n	$(1+x)^n$
complete tripartite graph $K_{l,m,n}$	$(1+lx)(1+mx)(1+nx)$
crossed prism graph	$3nx^2+2nx+1$
crown graph	$n(n-1)x^2+2nx+1$
cycle graph C_n	$\begin{cases} (1+x)^3 & \text{for } n=3 \\ 1+nx+nx^2 & \text{otherwise} \end{cases}$
empty graph \bar{K}_n	$nx+1$
folded cube graph	$\begin{cases} (x+1)^{2(n-1)} & \text{for } n=2, 3 \\ 1+2^{n-2}x(nx+2) & \text{otherwise} \end{cases}$
gear graph	$3nx^2+x(2n+1)x+1$
grid graph $P_m \times P_n$	$1+mnx+(2mn-m-n)x^2$
grid graph $P_l \times P_m \times P_n$	$1+lmnx+(3lmn-lm-ln-mn)x^2$
helm graph	$\begin{cases} (1+x)(1+6x+3x^2+x^3) & \text{for } n=3 \\ (1+x)(1+2nx+nx^2) & \text{otherwise} \end{cases}$
hypercube graph Q_n	$2^{n-1}x(2+nx)+1$
$m \times n$ -king graph	$\begin{cases} (1+x)^2(1+(n-1)x(2+x)) & \text{for } m=2 \\ (1+x)(1+(mn-1)x+3(m-1)(n-1)x^2+(m-1)(n-1)x^3) & \text{otherwise} \end{cases}$
$m \times n$ -knight graph	$1+mnx+2(2mn-3m-3n+4)x^2$
ladder graph $P_2 \square P_n$	$1-2x^2+nx(2+3x)$
ladder rung graph nP_2	$1+nx(2+x)$
Möbius ladder M_n	$1+nx(2+3x)$
path graph P_n	$(1+x)(1+(-1+n)x)$
prism graph Y_n	$\begin{cases} (2x+1)(x^2+4x+1) & \text{for } n=3 \\ 1+nx(2+3x) & \text{otherwise} \end{cases}$
star graph S_n	$(1+x)(1+(-1+n)x)$
sun graph	$nx(1+x)^2+(1+x)^n$
sunlet graph $C_n \odot K_1$	$1+2nx(1+x)$
transposition graph	$1+n!x+\frac{1}{4}n!n(n-1)x^2$
triangular grid graph	$\frac{1}{2}(1+x)(2+nx(3+n+2nx))$
web graph for $n > 3$	$1+3nx+4nx^2$
wheel graph W_n	$\begin{cases} (1+x)^4 & \text{for } n=4 \\ (1+x)(1+(-1+n)x(1+x)) & \text{otherwise} \end{cases}$

Chapter 1 Introduction

1.1 Background and Overview

Graph theory is a mathematical discipline that studies the properties and relationships of graphs, which are mathematical structures used to represent networks, relationships, and connectivity in various fields. A fundamental concept in graph theory is that of graph polynomials, which provide valuable insights into graph structures. This thesis aims to provide a comprehensive exploration of different types of graph polynomials and the methods for obtaining them. To appreciate graph polynomials significance, it is essential to delve into their historical background and understand how they came about. Graph theory can be traced back to the 18th century when the Swiss mathematician Leonhard Euler solved the famous "Seven Bridges Konigsberg" problem. Euler's work laid the foundation for graph theory by introducing the concept of a graph as a set of points connected by edges. In the city of Konigsberg there were seven bridges that had been built along the banks of the Pregel River. Euler showed how it was possible to cross all of these bridges only once. The field remained relatively dormant until the mid-20th century when the importance of graphs became apparent in various applications, including computer science, chemistry, and social sciences.

One of the earliest and most fundamental graph polynomials is the chromatic polynomial. One of the pioneering figures in this field was George David Birkhoff, an American mathematician who introduced the concept of the chromatic polynomial in 1912. Birkhoff's work laid the foundation for the study of graph polynomials. The chromatic polynomial, is defined for a graph G and represents the number of ways to color its vertices using k colors in such a way that no two adjacent vertices

have the same color. This polynomial is particularly important in graph coloring problems, which have applications in scheduling, map coloring, and register allocation in compiler design.

The idea of graph polynomials gained further traction with the introduction of the Tutte polynomials by W. T. Tutte in the 1940s. Tutte's polynomial, is a comprehensive tool that captures various aspects of a graph, including its connectivity, the number of spanning trees and more. It played a crucial role in development of graph theory and has applications in diverse fields, including electrical network analysis and statistical physics.

As graph theory continued to evolve, researchers developed the other graph polynomials, each tailored to address specific graph properties and problems. For example, the matching polynomial is used to count the number of matchings (independent edges) in a graph, making it relevant in combinatorial mathematics and network flow problems. The Jones polynomials, originating in knot theory, offers insights into the topological properties of knots and links.

Algebraic graph theory further expanded the realm of graph polynomials. The characteristic polynomial of a graph, derived from its adjacency matrix, provides information about the graph's spectrum and eigenvalues. This polynomial has applications in understanding the connectivity and structural properties of graphs.

In contemporary times, graph polynomials have become indispensable in numerous disciplines. In computer science, they play a pivotal role in solving optimization problems like finding maximum cliques, independent sets, and network flows. In the realm of network analysis, these polynomials help analyze complex systems such as social networks, transportation networks, and the internet, aiding in the identification of vulnerabilities and the study of resilience.

Graph polynomials also find applications in cryptography, where they assist in developing secure communication protocols. In scientific modeling, they are used to represent molecular structures, crystal lattices, and other complex systems in chemistry and physics. Furthermore, they are integral to graph-based machine learning models that are increasingly used in data analysis and artificial intelligence.

Graph polynomials have a rich history and have evolved into a fundamental mathematical tool with a wide range of applications. Their development over the years has been driven by the need to understand, analyze, and solve problems related to graphs in various scientific, engineering, and real-world contexts. Graph polynomials continue to be a vital area of research and a key component of modern mathematics and its applications.

1.2 Definitions

1. Graph : A graph is a mathematical structure consisting of a set of vertices (nodes) and a set of edges (connections) that defines pairwise relationships between the vertices.

2. Polynomial: A polynomial is a mathematical expression consisting of variables and coefficients, combined using addition, subtraction, and multiplication operations. In the context of graph theory, polynomial expressions often represent certain graph properties.

3. Graph Polynomial : A graph polynomial is a polynomial associated with a graph, where the coefficients and variables have specific meanings which are used to capture and quantify the graph's structure or combinatorial properties.

4. Chromatic Polynomial : The chromatic polynomial of a graph G , denoted as $P(G, \lambda)$, is a polynomial that counts the number of ways to color the vertices of G with λ colors such that no two adjacent vertices share the same color.

5. Characteristic Polynomial : The characteristic polynomial of a graph G , denoted as $\chi(G, \lambda)$, is a polynomial used in the study of graph spectra. It provides information about the eigenvalues of the adjacency matrix of G .

6. Tutte Polynomial: The Tutte polynomial of a graph G , denoted as $T(G, x, y)$, is a polynomial that encodes a variety of graph properties.

7. Independence polynomial : Counts the number of independent sets of vertices in a graph of different sizes.

8. Matching polynomial : Counts the number of matchings of different sizes in a

graph. A matching is a set of edge, no two of which share a vertex.

9. **Polynomial Coefficients:** Coefficients of a graph polynomial often have combinatorial significance and relate to structural properties of the graph. 10. **Graph Isomorphism:** Graph polynomial can be used to distinguish non-isomorphic graphs with the same polynomial.

1.3 Introduction to Graph Polynomials

Graph polynomials are mathematical tools used in the study of graphs, which are mathematical structures consisting of nodes (vertices) and edges that connect those nodes. These polynomials assign a numerical value to a graph, capturing various graph properties, and they have important applications in a variety of fields, including computer science, chemistry, physics, and social network analysis.

Graph polynomials are important to learn about in contemporary times for several reasons:

1. **Network Analysis:** In the age of information and communication networks, graph polynomials play a crucial role in the analysis of social networks, transportation networks, and the internet. They help understand the connectivity, vulnerabilities, and resilience of these complex systems.
2. **Cryptography:** Graph polynomials are used in cryptographic protocols and coding theory for error detection and correction. They are essential in designing secure communication systems.
3. **Scientific Modeling:** In various scientific disciplines, including biology, chemistry, and physics, graph polynomials are used to model and analyze molecular structures, crystal lattices, and other complex systems.
4. **Computer Science:** In computer science and algorithms, graph polynomials are employed to solve optimization problems, such as finding the maximum clique in a graph or the maximum independent set.

5. **Machine Learning:** Graph-based machine learning models, such as graph neural networks, are becoming increasingly important in various applications. Understanding graph polynomials can provide insights into these models' behavior.
6. **Social Sciences:** Graph theory and its associated polynomials are used to analyze social networks, study information diffusion, and understand the dynamics of online communities.

Graph polynomials are a fundamental and versatile tool in mathematics and various scientific disciplines. They help us analyze and model complex systems, solve optimization problems, and make sense of the interconnected world we live in today. Learning about graph polynomials is essential for anyone interested in understanding and solving real-world problems in a networked and data-driven society.

1.4 Importance of Graph Theory

Graph theory has immense importance in both theoretical and practical contexts. Its versatility and applicability make it an indispensable tool for understanding and solving complex problems.

The significance of graph theory in various domains, emphasize its impact in the real world.

1. **Computer Science and Networks:** Graph theory is the foundation of computer science, powering algorithms, data structures, and network modeling. In computer networks, it helps design efficient routing protocols, ensuring data travels optimally across the internet. Social networks, a fundamental part of the digital age, are analyzed using graph theory to understand information flow, identify influencers, and improve user experience. Graph algorithms are integral to solving problems like finding the shortest path, network connectivity, and spanning tree construction.
2. **Transportation and Logistics:** Graphs are employed in modeling transporta-

tion systems, optimizing routes for delivery trucks, and scheduling public transportation. The traveling salesman problem (TSP), a classic optimization challenge, is tackled using graph theory to find the most efficient route for visiting a set of locations. Such applications improve resource allocation, reduce costs, and enhance overall efficiency in transportation and logistics.

3. **Social Sciences and Communication:** Graph theory is instrumental in analyzing social networks and communication structures. By examining connections between individuals, researchers gain insights into information diffusion, influence propagation, and community detection. These findings are applicable in sociology, marketing, and even national security, where identifying patterns in communication is vital.
4. **Biology and Medicine:** In biology, graph theory aids in modeling genetic networks, protein-protein interactions, and metabolic pathways. Understanding these complex biological systems is crucial for drug discovery, disease diagnosis, and personalized medicine. Graphs also represent brain connectivity networks, helping neuroscientists study brain function and connectivity, leading to advances in neuroscience and mental health research.
5. **Operations Research and Optimization:** Graph theory provides essential tools for solving optimization problems. Whether it's optimizing the allocation of resources in supply chain management or finding the best configuration in manufacturing, graph-based algorithms help organizations make informed decisions, improve efficiency, and reduce costs.
6. **Engineering and Circuit Design:** Electrical engineers use graph theory to design circuits, analyze network flow, and optimize power distribution. By representing circuits as graphs, engineers can identify issues, optimize performance, and ensure the reliability of electrical systems.
7. **Geography and GIS:** Geographic Information Systems (GIS) heavily rely on graph theory for map routing and spatial analysis. It assists in finding the

shortest path between two locations, planning transportation networks, and understanding spatial relationships, making it indispensable for urban planning, disaster management, and environmental science.

8. **Linguistics and Natural Language Processing (NLP):** In linguistics, graphs help analyze syntactic and semantic relationships in language. In NLP, they are used for text summarization, sentiment analysis, and information retrieval. Understanding language structures and connections enhances machine translation, chatbots, and information retrieval systems.
9. **Finance and Economics:** Graph theory is applied in risk assessment, portfolio optimization, and fraud detections within the financed industry. By modeling financial relationships and transaction flows, it aids in identifying irregularities and optimizing investment strategies.
10. **Cybersecurity:** In the realm of cybersecurity, graph theory assists in detecting network intrusions, identifying patterns of malicious activity and securing critical systems. It is used to build attack graphs, visualizing potential vulnerabilities and attack paths.

Graph theory is not just a theoretical branch of mathematics; it's a practical and interdisciplinary tool that underpins modern society. Its ability to model, analyze, and solve complex problems across various domains, from computer science, to biology, economics, and beyond, demonstrates its unparalleled importance. Graph theory continues to drive innovation, improve efficiency, and provide invaluable insights into the interconnected world we live in, making it a cornerstone of contemporary mathematics and science.

Chapter 2 Different Types of Graph Polynomials

2.1 Chromatic and Coloring Polynomials

2.1.1 The Chromatic Polynomial

The *chromatic polynomial*, denoted as $P(G; x)$, is a polynomial associated with a given graph G that provides information about the number of ways to color its vertices with a specified number of colors. It is defined recursively based on the number of colors used. They are used to address problems related to scheduling and optimization in various domains. They help determine a graph's chromatic number.

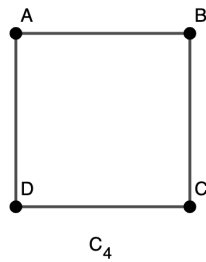
Understanding chromatic polynomials aids in assessing the structural properties of graphs and their coloring behavior.

- 1. **Monotonicity:** The chromatic polynomial is a monotonically decreasing function. In other words, as the number of colors (k) increases, the number of valid colorings ($P(G, k)$) decreases.
- 2. **Evaluation:** Calculating $P(G, k)$ for a graph is a complex task. Algorithms like the deletion-contraction recurrence or polynomial interpolation can be employed to determine the chromatic polynomial.
- **Zero Property:** The chromatic polynomial evaluates to zero for k less than the chromatic number of the graph. This means that for $k < \chi(G)$, there are no valid colorings of the graph.
- **Chromatic Roots:** The chromatic polynomial can be factored into linear terms as $(k - \lambda_1)(k - \lambda_2) \dots (k - \lambda_n)$, where λ_i represents a chromatic root. The chromatic roots provide insights into the colorability of the graph.

The study of chromatic polynomials continues to be a vibrant area of research, furthering our knowledge of graph theory and its applications. As we explore more complex networks and structures, the significance of chromatic polynomials in graph theory becomes increasingly pronounced.

By examining the coefficients and roots of these polynomials, mathematicians and computer scientists gain valuable insights into the graph's colorability, planarity, and connectivity.

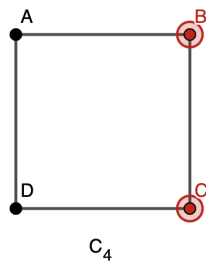
Example



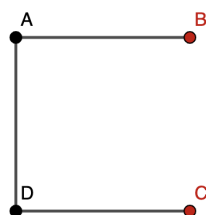
The graph that we will be using to find the chromatic polynomial will be a cycle graph with four vertices. C_4

Step 1: Pick two vertices that are adjacent to each other.

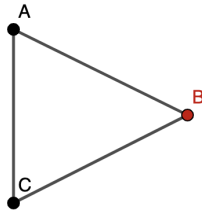
In this example, I will be showing it with the vertices B and C.



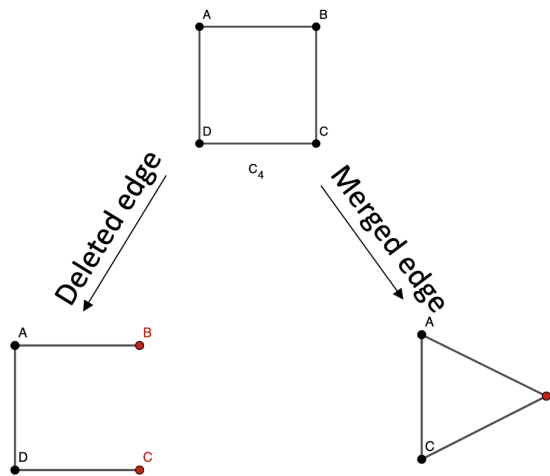
Step 2: We will build two graphs based off of the decision of the chosen vertices. The first graph will show where we deleted the edge that were adjacent to the vertices B and C.



Now that we have our deleted edge we will move on and merge the 2 vertices together.



Now that we have our graphs this is what we have:



Looking at the two graphs that we have now, we can see that graph to our left is a path graph and the graph that we merged is a complete graph.

The formula to finding the chromatic number for a path graph is as follows:

$$P_x(P_n) = x(x - 1)^{n-1}, \text{ where } n \text{ is the number of vertices.}$$

We have a path with 4 vertices so the polynomial for the deleted edge graph is:

$$P_x(P_4) = x(x - 1)^3$$

We know for the complete graph for the chromatic polynomial the formula is the fallen factorial.

So, for the merged edge graph, the chromatic polynomial we know for a complete graph with 3 vertices would be:

$$P_x(K_3) = x(x - 1)(x - 2)$$

Now that we have both of our polynomials we can now combine them together for deletion-contraction method.

So, we now have:

$$\begin{aligned} & x(x-1)^3 - (x(x-1)(x-2)) \\ & x^4 - 3x^3 + 3x^2 - x - x^3 + 3x^2 - 2x \\ & x^4 - 4x^3 + 6x^2 - 3x \end{aligned}$$

So our chromatic polynomial for this cycle graph is:

$$P_x(C_4) = x^4 - 4x^3 + 6x^2 - 3x$$

2.1.2 Clique Polynomial

A clique is a subset of vertices in a graph such that every pair of vertices in a subset is connected by an edge. The concept of cliques is essential for understanding more complex structures in graphs. The clique polynomial is a polynomial associated with a graph that provides information about the graph's clique structure. Given a graph G , the clique polynomial is typically denoted as $\omega(G, x)$. It is defined as the sum of a series of terms, each representing a clique of the graph G . The expression for the clique polynomial is as follows:

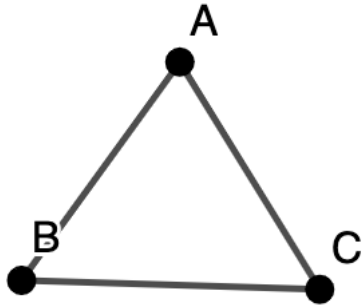
$$\omega(G, x) = \sum_{k=0}^M a_k x^k \tag{2.1}$$

While clique polynomials offer valuable insights into graph structure, they are not without their challenges. Analyzing the clique polynomials involve applying them to more complex graph structures, such as hypergraphs, where edges can connect to more than two vertices. Additionally, combining clique polynomials with machine learning techniques may open new avenues for graph analysis and prediction in various domains. Understanding the clique polynomial of a graph provides valuable insights into the underlying structure and relationships within the graph.

Example:

Step 1: Understand the graph

Start by understanding the simple graph for which you want to find the clique polynomial. A simple graph consists of a set of vertices and a set of edges that connect these vertices.



Step 2: Identify Cliques

A clique is a subset of vertices in the graph where every pair of vertices is adjacent (connected by an edge). You want to find all the possible cliques in the graph.

Clique polynomials are built from the number of cliques in an undirected simple graph.

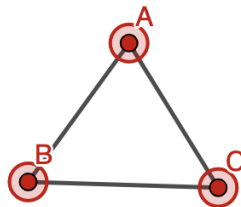
The clique polynomials of an undirected graph G is defined as:

$a_0 + a_1x^1 + a_2x^2 + \dots a_nx^n$ where a_n is the number of n -cliques in G .

We know that a_0 , or the number of 0 cliques, is always 1 for any graph, and that this polynomial terminates with the term corresponding to the largest clique(s) in the graph, an n -clique.

So, we have $a_0 = 1$

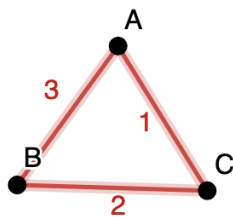
To find the number of 1-cliques we are looking to see how many vertices are showing in the graph.



By looking at the graph we see that there are 3 vertices : A, B, and C. This means that there are 3 1-cliques in this graph.

So, we have $a_1 = 3$

To find the number of 2 - cliques we are looking to see how many edges are showing in the graph.

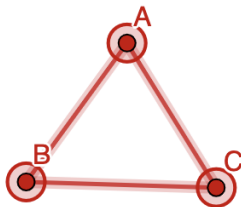


We have 3 edges in this graph (A-C), (B-C), and (A-B).

So, we have $a_2 = 3$

Last, we know in this graph that our M is equal to 3.

By observing our graph, we can see, that our entire graph K_3 , forms a clique.



We only have one, so $a_3 = 1$

We are now able to put our clique polynomial together.

$$a_0 = 1$$

$$a_1 = 3$$

$$a_2 = 3$$

$$a_3 = 1$$

The clique polynomial of this graph would be:

$$\omega(G, x) = a_0 + a_1x^1 + a_2x^2 + a_3x^3 \tag{2.2}$$

$$\omega(G, x) = 1 + 3x^1 + 3x^2 + 1x^3 \tag{2.3}$$

Easiest way to interpret this polynomial is knowing that the coefficients in this polynomial represents the cliques of each size of the graph as well you interpreting the number of vertices, edges and triangles in the graph.

# of vertices in the graph	# of edges in the graph	# of triangles in the graph
-------------------------------------	----------------------------------	--------------------------------------

$$\omega(G, x) = 1 + 3x^1 + 3x^2 + 1x^3$$

Keep in mind that finding the clique polynomial can become more complex for larger or more intricate graphs, but the process remains the same: identify cliques of different sizes and construct the polynomial accordingly.

2.1.3 Rainbow Connection Polynomial

The *Rainbow Connection Polynomial*, denoted as $RC(G, k)$, is a polynomial that encodes information about the rainbow connection number of a graph G concerning the number of colors used, k . It has computed as a generating function that counts the number of graphs with rainbow connection number k . The polynomial $RC(G, k)$ is defined as follows:

$$RC(G, k) = \sum_{G' \subset G} [C(G')]k, \tag{2.4}$$

where the sum is taken over all subgraphs G' of G , and $C(G')$ is the chromatic polynomial of G' , which represents the number of proper edge colorings of G' using k colors.

It provides a quantitative measure of connectivity of a graph. While traditional connectivity metrics focus on paths or components, the rainbow connection number and polynomial offer a unique perspective by considering the diversity of paths between vertex pairs. They offer ways to quantify and measure graph connectivity, which is essential for solving real-world problems that involve networks and connections. The Rainbow Connection Polynomial is interconnected with other graph polynomials, such as the chromatic polynomial and the Tutte polynomial.

2.1.4 Heterochromatic Polynomial

The *heterochromatic polynomial* is a concept closely related to the chromatic polynomial but with a unique twist. In the context of a heterochromatic polynomial, we are not only interested in coloring the vertices of a graph with colors but also specifying colors for certain vertices in advance.

Let's delve deeper into this concept. Consider a graph G , where we have two sets of vertices: set A and set B . Set A contains certain vertices whose colors are predetermined, while set B includes the remaining vertices whose colors must be determined according to the proper coloring rule.

The heterochromatic polynomial, denoted as $H(G, x, y)$, is defined as a polynomial in two variables, x and y . This polynomial represents the number of ways to properly color the vertices in set B with x colors while ensuring that the vertices in set A are colored with y distinct colors. The term "heterochromatic" arises from this dual-coloring property, where the colors for sets A and B are different.

Heterochromatic colorings are more restrictive than proper colorings, as they require not only adjacent vertices but also non-adjacent vertices to have different colors. It can be defined by using a recursive formula. Let G be a graph with n vertices, and v be any vertex in G . The heterochromatic polynomial $H(G, x, y)$ can be computed as follows:

$$H(G, x, y) = \sum (\text{over all proper 2-colorings of the graph } G) (x^k * y^{(n-k)}) \quad (2.5)$$

Where:

- k is the number of vertices colored with color x .
- n is the total number of vertices in the graph ($|V|$)
- The sum is taken over all possible 2-colorings of the graph.

The heterochromatic polynomial exhibits several noteworthy properties that make it significant as a graph polynomial:

- Symmetry: $H(G, x, y)$ is symmetric, which implies that the order of x and y in the colorings doesn't affect the result.
- Initial Values: For a graph with no edges, $H(G, x, y) = x^y$, reflecting the number of possible colorings using x and y colors.
- Degree: The degree of $H(G, x, y)$ is at most n , where n is the number of vertices in the graph G .
- Chromatic Polynomial: When $x = y$, the heterochromatic polynomial reduces to the chromatic polynomial of the graph, which counts the number of proper colorings using x colors. As research continues to unfold, the heterochromatic polynomial remains a key tool for both theoretical graph theory and practical problem-solving in the real world. Its significance as a graph polynomial lies in its ability to illuminate complex relationships within graphs and provide solutions to a diverse range of combinatorial challenges.

2.2 Structural Polynomials

2.2.1 Tutte Polynomial

Tutte polynomials are named after the Canadian mathematician W. T. Tutte. These polynomials are used to study and analyze various properties of graphs, providing insights into their structure, connectivity, and other combinatorial characteristics.

The Tutte polynomial of a graph G , denoted as $T(G; x, y)$, is a bivariate polynomial defined in terms of two variables, x and y . It encodes information about the graph's different combinatorial aspects, such as the number of spanning trees, the number of independent sets, and the chromatic polynomials. They have wide applications in graph theory, particularly in the study of planar graphs, matroids, and electrical networks. They are used to solve problems related to network reliability, graph coloring, and counting spanning trees. They have applications in statistical physics, where they are used to analyze phase transitions in physical systems. The Tutte polynomial is typically defined using a recursive approach. Given a graph G ,

let $|G|$ represent the number of vertices in the graph. The base case of the recursion is as follows:

- i. $T(G; x, y) = 1$ if G is the empty graph (no vertices or edges)
- ii. $T(G; x, y) = xT(G - e; x, y)$ if G can be obtained from $G - e$ by contracting an edge e .
- iii. $T(G; x, y) = yT(G - e; x, y)$ if G can be obtained from $G - e$ by deleting an edge e .
- iv. $T(G; x, y) = T(G - e_1 - e_2; x, y)$ if G can be obtained from $G - e_1$ and $G - e_2$ by gluing them together along a common vertex.

Where $G - e$ represents the graph obtained by deleting edge e , and $G - e_1 - e_2$ represents the graph obtained by deleting edges e_1 and e_2 . One of the primary uses of the Tutte polynomial is in enumerating subgraphs with specific properties. This enumeration is valuable in combinatorics and network analysis, as it allows us to understand the prevalence of various structural components within a graph. Its duality with the chromatic polynomial and its connections to matroid theory further illustrate its importance in the broader landscape of combinatorial mathematics. As we continue to explore the world of graphs and networks, the Tutte polynomial remains an indispensable asset, allowing us to unlock the secrets hidden within these interconnected structures.

The Tutte polynomial's versatility lies in its ability to unify several important combinatorial problems under a single mathematical framework. Researchers use it to gain insights into the relationships between these problems and to develop efficient algorithms for solving them.

More recently, the study of graph polynomials has expanded to include other types, such as the Jones polynomial in knot theory, the matching polynomial in combinatorial mathematics, and the characteristic polynomial in algebraic graph theory.

2.2.2 Flow Polynomial

Flow polynomials are used in graph theory to analyze and understand the flow of objects or information within networks. These networks are represented as graphs, with nodes representing entities and edges representing connections or pathways between them. Flow polynomials provide a concise way to express and compute various properties related to flow in these networks.

Flow polynomials have wide applications in various fields, including transportation planning, communication networks, and resource allocation. They can be used to find optimal flow distributions, assess network capacity, and solve various optimization problems.

Overall, flow polynomials play a crucial role in the analysis and optimization of network flows, making them an indispensable tool in fields where understanding the movement or distribution of resources or information is essential.

2.2.3 Matching Polynomial

Matching polynomials are used to study the properties and characteristics of matchings within graphs. In graph theory, a matching is a set of edges in a graph where no two edges share a common vertex. The matching polynomial, denoted as $M(G, x)$, is a polynomial associated with a given graph G , and it provides valuable insights into the graph's matching structure.

The matching polynomial is defined as the sum of the matching numbers for all possible cardinalities of matchings in the graph. They have several important applications in various domains including computer science, biology, and network design. They can be used to solve problems related to finding maximum matching, which are crucial in tasks like optimizing network flows, scheduling and resource allocation. They are also used in the analysis of structural properties of graphs, such as determination of bipartite structure and the assessment of graph connectivity.

Matching polynomials can be used to study the combinatorial properties of graphs, allowing researchers to make conclusions about the graph's structure and behavior

based on its polynomial representation. By examining roots, coefficients, and other properties of the matching polynomial, mathematicians and computer scientists gain a deeper understanding of the graph's matching characteristics.

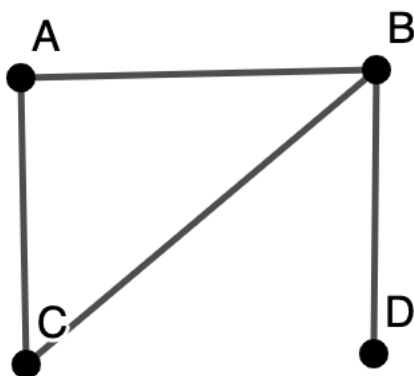
2.2.4 Independence Polynomials

Independence Polynomial represents the number of independent sets of size k in G , where k ranges from 0 to the number of vertices in the graph. It is used to study the structure and independence properties of graphs.

An independent set is a set of vertices in which no two vertices are adjacent. Finding the independence polynomial involves counting the number of independent sets of different sizes in the graph.

Example: Let's consider a simple graph with 4 vertices (A, B, C and D) and the following edges:

1. A-B
2. A-C
3. B-C
4. B-D



Step 1: Create an adjacency matrix First, create an adjacency matrix to represent the graph. The adjacency matrix is a square matrix where the entry $A[i][j]$ is 1 if there is an edge between the vertices i and j , and 0 otherwise. In this example, the adjacency matrix would look like this:

	A	B	C	D
A	0	1	1	0
B	1	0	1	1
C	1	1	0	0
D	0	1	0	0

Step 2: Calculate the independence polynomial To find the independence polynomial, you will calculate the number of independent sets of different sizes in the graph.

- Independent sets of size 0 (the empty set): There is only one way to choose an empty set, which is by not selecting any vertices. So, the coefficient of x^0 is 1.
- Independent sets of size 1: There are four vertices in the graph (A, B, C, and D), and each can form an independent set by itself. So, the coefficient of x^1 is 4.
- Independent sets of size 2: You need to count the number of 2-vertex independent sets. In our graph, there is one pair (A, D) that can form an independent set. So, the coefficient of x^2 is 1.
- Independent sets of size 3: There are no 3 - vertex independent sets in this graph. So, the coefficient of x^3 is 0.
- Independent sets of size 4: There is only one independent set of size 4 in this graph, which is the set A, B, C, D. So, the coefficient of x^4 is 1.

Now we construct the independence polynomial by summing up the terms:

$$\text{Independence polynomial} = 1 + 4x + x^2 + 0x^3 + x^4$$

So, the independence polynomial for this graph is:

$$P(x) = 1 + 4x + x^2 + x^4$$

This polynomial encodes information about the number of independent sets of different sizes in the graph. In more complex graphs, finding this polynomial can be

more involved, but the process remains the same - count the number of independent sets for each size and construct the polynomial accordingly.

2.3 Matrix-Based Polynomials

2.3.1 Adjacency Matrix (Characteristic Polynomial)

The characteristic polynomial of a graph, often denoted as $\det(G - \lambda I)$, where G is the graph's adjacency matrix, I is the identity matrix, and λ is a variable, is a polynomial that captures information about the graph's eigenvalues. The roots of the characteristic polynomial are the eigenvalues of the graph. Studying the eigenvalues of a graph provides insights into its spectral properties, including connectivity, expansion, and graph partitioning. The spectral graph theory heavily relies on characteristic polynomials to analyze various aspects of graphs, such as random walks, Laplacian matrices, and clustering algorithms.

Characteristic polynomials are versatile tools that bridge the gap between algebraic structures and geometric or structural properties. They are essential in various applications, including physics, computer science, and engineering, where they help analyze linear transformations, study network properties, and solve differential equations. By examining the roots and coefficients of characteristic polynomials, mathematicians and scientists gain valuable insights into the systems they represent making them a fundamental concept in the study of matrices, graphs, and linear systems.

2.3.2 Kirchhoff Polynomial

The Kirchhoff polynomial is a special graph polynomial named after the renowned German physicist Gustav Kirchhoff, who made significant contributions to various fields. The polynomial is formally defined as the determinant of a matrix derived from the Laplacian matrix of a graph. The Kirchhoff polynomial, denoted $K(G, x)$, is obtained by taking the determinant of the Kirchhoff matrix $K - xI$, where I is the identity matrix of the same size as K .

$$K(G, x) = \det(K - xI)$$

In this equation, x is a variable that can take on different values, and G represents the graph under consideration.

The Kirchhoff polynomial is not an isolated concept. It is connected to other graph polynomials. It can be expressed in terms of the characteristic polynomials and chromatic polynomial of a graph. It has applications in counting spanning trees, analyzing graph connectivity, understanding vertex and degrees and edge cuts.

2.3.3 Resistance-Distance Polynomial

The resistance distance polynomial is a relatively new addition to the family of graph polynomials, offering valuable insights into electrical networks, network connectivity, and structural properties of graphs. Resistance distance between two vertices u and v in a graph G is a measure of how difficult it is for electrical current to flow from u to v . It is defined as the difference in energy (voltage) when a unit current is injected at u and extracted at v . The resistance distance $D(u, v)$ can be expressed as:

$$D(u, v) = \frac{1}{I_{uv}} - \frac{1}{I_{vv}}$$

Where I_{uv} is the current between u and v and when a unit voltage is applied at u , and I_{vv} is the current between v and v when a unit voltage is applied at u . The resistance distance polynomial, denoted as $R(G, x)$, is a graph polynomial associated with the resistance distances in a graph. It is a polynomial of a real variable x and provides valuable information about the graph's structure and connectivity. The $R(G, x)$ polynomial is defined as follows:

$$R(G, x) = \sum_{u,v} D(u, v)x^{d(u,v)} \quad (2.6)$$

Where the sum is taken over all pairs of vertices u and v , and $d(u, v)$ represents the distance between u and v in the graph G . The resistance distance polynomial encapsulates the entire spectrum of resistance distances in a graph, offering a unique perspective on graph properties. Its significance lies in its ability to aid in graph analysis, network design, structural property determination, and various applications

in science and engineering. As graph theory continues to find applications in a wide range of fields, the resistance distance polynomial remains a valuable tool for understanding the hidden structures and properties of graphs.

2.4 Directed Graph Polynomials

2.4.1 Seidel-Entringer-Arnold Polynomial

The Seidel polynomial is named after the German mathematician Richard Seidel, who introduced it in the 1970s. It is defined for a graph G and a positive integer k as the Seidel matrix, $S(G, k)$, which represents a modified adjacency matrix of G . Specifically, $S(G, k)$ is obtained by replacing the diagonal entries of the adjacency matrix $A(G)$ with the degree of each vertex and then raising the resulting matrix to the power of k . The Seidel polynomial is defined as:

$$P(G, x) = \sum_{k=0}^n (-1)^k \text{tr}(S(G, k)) x^k \quad (2.7)$$

Where $P(G, x)$ is the Seidel polynomial of the graph G , $\text{tr}(S(G, k))$ is the trace of the Seidel matrix $S(G, k)$, and n is the number of vertices in G . The Seidel polynomial was further extended by Richard Entringer, who introduced the concept of Entringer matrices. These matrices are closely related to Seidel matrices and provide another way to represent graph polynomials.

The Entringer polynomial is defined as:

$$Q(G, x) = \sum_{k=0}^n \text{tr}(A^k) x^k,$$

Where $Q(G, x)$ is the Entringer polynomial of the graph G , $\text{tr}(A^k)$ is the trace of the k th power of the adjacency matrix $A(G)$, and n is the number of vertices in G .

The Arnold polynomial, named after the mathematician Viatcheslav Arnold. Arnold's work on graph polynomials provides a unique perspective on the properties of graphs.

The Arnold polynomial is defined as:

$$R(G, x) = \sum_{k=0}^n \text{tr}(B^k) x^k,$$

Where $R(G, x)$ is the Arnold polynomial of the graph G , $\text{tr}(B^k)$ is the trace of

the k th power of the biadjacency matrix $B(G)$, and n is the number of vertices in G .

The Seidel-Entringer-Arnold Polynomial is known as the SEA polynomial. It represents a unique way of capturing graph properties by blending the Seidel, Entringer, and Arnold perspectives into a single polynomial.

The SEA polynomial is defined as:

$$S(G, x, y, z) = P(G, x) + Q(G, y) + R(G, z).$$

Here, $S(G, x, y, z)$ is the Seidel-Entringer-Arnold polynomial, $P(G, x)$ is the Seidel polynomial, $Q(G, y)$ is the Entringer polynomial, and $R(G, z)$ is the Arnold polynomial. The SEA polynomial can be used to develop algorithms for solving graph-related problems efficiently.

2.4.2 Flow-Path Polynomial

The primary goal of flow-path polynomials is to quantify the number of ways in which flow can be established between two specific nodes within the network. These paths represent various routes or channels that the flow can take, and the polynomial captures this information in a compact and analytical form.

In a directed graph which consists of nodes and direct edges, each edge is assigned a capacity, indicating the maximum amount of flow it can carry. A flow-path polynomial $F(x)$ for a graph F is a polynomial that counts the number of ways to send a flow from a source node(s) to a sink node (t) subject to the capacity constraints, while ensuring that no edge carries more flow than its capacity.

The flow-path polynomial $F(x)$ is defined as:

$$F(x) = \sum_f x^f$$

Where:

- f represents a feasible flow in the graph G .
- x^f represents the flow value associated with the feasible flow f .

This polynomial captures all possible ways flow can be routed from the source to sink while respecting the capacity constraints. It's important to note that each term in the polynomial $F(x)$ corresponds to a specific feasible flow configuration in the

network. Flow-path polynomials are a specialized type of graph polynomial, focusing on flow-related properties within a directed graph.

2.5 Geometric and Spectral Polynomials

2.5.1 Laplacian Spectral Density Polynomial

2.5.2 Q- Flow Polynomial

The Q-Flow Polynomial is a graph polynomial that focuses on the flow-related properties of a graph. It was introduced as a tool to study flows in networks and is particularly useful in the context of electrical networks and flow optimization problems. The Q-Flow Polynomial is derived from the Kirchhoff matrix of a graph, which represents the relationships between the flow of electrical currents and the voltage differences in a network. 1. The Kirchhoff matrix, also known as the Laplacian matrix, is a square matrix derived from a graph. It is defined as follows for an undirected graph G with n vertices:

$$L(G) = D(G) - A(G)$$

Where:

- $L(G)$ is the Laplacian matrix.
- $D(G)$ is the diagonal matrix of vertex degrees.
- $A(G)$ is the adjacency matrix

The Q-Flow Polynomial, denoted as $Q(G, t)$, is derived from the Kirchhoff matrix. It is defined as the determinant of the matrix $(tI - L(G))$, where I is the identity matrix and t is a variable:

$$Q(G, t) = \det(tI - L(G))$$

By evaluating $Q(G, t)$ for different values of t , researchers can gain insights into how flows propagate and distribute within the graph.

Q-Flow Polynomials complement other graph polynomials like the characteristic polynomial, chromatic polynomial and flow polynomials. Q-Flow Polynomials provide

insights into flow-related properties. As graph theory continues to find applications in diverse domains, the significance of Q-Flow polynomials in understand network work becomes increasingly evident.

2.5.3 Theta Graph Polynomial

Theta graphs are a special class of graphs that play a pivotal role in graph theory. A theta graph is formed by connecting three vertices (called the 'end vertices') to a common fourth vertex (the 'center vertex') using three edges, creating a structure that resembles the Greek letter theta (Θ). This simple definition masks the complexity and significant of theta graphs when it comes to graph theory.

Two prominent theta graph polynomials are the theta Polynomial ($T_G(x)$) and the extended theta polynomial ($T_G^*(x)$).

1. Theta Polynomial:

$(T_G(x)) = x^2 * (x - 1) * N_\theta(G)(x)$, where $N_\theta(x)$ is the characteristic polynomials of the subgraph formed by the end vertices. This polynomial represents the number of closed walks of different lengths in the theta graph.

2. Extended Theta Polynomial: ($T_G^*(x)$):

The extended theta polynomial of a graph G is a more general form, defined as:

$$(T_G^*(x)) = x^3 * (x - 1) * N_\theta(G)(x).$$

This polynomial provides additional information about the graph, including its eigenvalues and more complex structural properties. Theta graph polynomials provide a concise and powerful representation of the structural properties of graphs. They can be used to determine various parameters. They are also essential in graph isomorphism testing. By comparing the theta polynomials of two graphs, one can determine if they are isomorphic. By capturing the essence of theta graphs in a polynomial form, these mathematical expressions serve as a bridge between abstract graph theory and practical applications.

Chapter 3 Applications of Graph Polynomials

3.1 Network Analysis

Networks, including social networks, transportation networks, and communication networks, can be modeled as graphs. Graph polynomials play a crucial role in network analysis.

- **Connectivity Polynomials:** The polynomials, like the Kirchhoff polynomial, provide insights into the connectivity properties of a graph.
- **Reliability Analysis:** Graph polynomials are used to assess the reliability of network systems, such as communication networks, by quantifying the probability of successful communication in the presence of network component failures.
Connectivity: The characteristic polynomial can determine whether a graph is connected by examining the existence of a zero eigenvalue in the associated adjacency matrix.

Community Detection: Spectral analysis using characteristic polynomials helps identify communities or clusters within large networks. Resilience Analysis: Graph polynomials can assess the robustness of networks against node or edge failures.

3.2 Coding Theory

- **Matriod Polynomials:** Graph polynomials are used in coding theory to design error-correcting codes, optimize communication systems, and study the structural properties of codes.
- **Connectivity and Rank Metric Codes:** In the study of rank metric codes,

graph polynomials help analyze the minimum distance properties of codes, which are crucial in error correction.

3.3 Social Sciences

In social sciences, graph theory and graph polynomials are applied to model and analyze social networks and relationships:

- **Social Network Analysis:**

Chromatic polynomials are used to identify the number of overlapping communities within a social network. Influence Propagation: Graph polynomials assist in studying the spread of information or influence within social networks.

- **Opinion Dynamics:** Characteristic polynomials can provide insights into the stability of opinion dynamics in a society.

3.4 Network Security

- **Anomaly Detection:** By employing graph polynomials, network administrators can detect anomalies and malicious activities in network traffic, identifying patterns that deviate from the norm.

3.5 Computer Science

Graph theory has extensive applications in computer science, and graph polynomials are no exception:

- **Algorithm Analysis:** Graph polynomials help analyze the time complexity of graph algorithms and data structures, providing insights into computation efficiency.

- **Isomorphism Testing:** Characteristic polynomials can be used to compare graph structures for isomorphism, which is essential in graph database searches and pattern recognition.

- **Network Routing:** Chromatic polynomials assist in solving network routing problems, where data packets need to traverse a graph with certain constraints.

3.6 Image Processing

- **Graph Matching:** Graph polynomials facilitate graph matching and pattern recognition, where the similarity between graphs is assessed using algebraic invariants.
- **Image Segmentation:** Graph polynomials can help partition an image into distinct regions based on pixel connectivity.
- **Feature Extraction:** Characteristic polynomials are used to extract structural features from images for object recognition and classification.

3.7 Computational Biology

- **Protein-Protein Interaction Networks:** Graph polynomials assist in analyzing complex biological networks, such as protein-protein interaction networks, to uncover functional relationships and regulatory mechanisms.
- **Phylogenetics:** In the study of evolutionary relationships among species, phylogenetic trees can be analyzed using graph polynomials to infer common ancestry and genetic evolution.

Spectral Connections: The Laplacian polynomial is closely related to the spectral properties of a graph. It helps analyze eigenvalues, and the multiplicity of the eigenvalue zero corresponds to the number of connected components in the graph. Spectral techniques can be used in network reliability analysis as well. Characteristic polynomials are also closely linked through eigenvalues of the graph. Spectral analysis bridges the gap between these two type of polynomials. **Network Flow and Reliability:** Flow polynomials, especially Tutte matrix polynomials, are related to the

concept of spanning trees, which is crucial in network reliability analysis. The number of spanning trees in a graph is often used to assess network reliability. Reliability polynomials are more specialized in context of network analysis. Flow polynomials help analyze network flows and connectivity, while reliability polynomials focus on system reliability under failure conditions. Combinatorial Optimization: Both flow polynomials and Kirchhoff polynomials have applications in combinatorial optimization problems, which involve finding optimal solutions to discrete problems. They are used to optimize network structures and solve various combinatorial problems. Network Design: All three types of polynomials play a role in network design. Laplacian polynomials help identify influential nodes, flow polynomials aid in optimizing network flows, and reliability polynomials guide decisions about network redundancy and fault tolerance.

These different graph polynomials serve distinct purposes and are used to address various aspects of graphs and networks. While there may be some overlap and connections among them, each polynomial type emphasizes specific graph properties and analytical techniques. Understanding these differences is essential for applying the appropriate polynomial in a given context or problem-solving scenario.

Chapter 4 Real-World Cases

4.1 Case Study 1: Facebook's use of graph polynomials for friend recommendation

Facebook, as one of the world's largest social networking platforms, is constantly seeking ways to improve user experience and engagement. A key component of this is friend recommendations. In the context of Facebook, friend recommendations are made by analyzing the social graph - a mathematical representation of the connections between users. A social graph is essentially a network that represents relationships between users. In Facebook's case, each user is a node, and their connections with other users (i.e, friendships) are represented as edges. This results in a massive, interconnected graph with millions of nodes and edges, making it a complex mathematical object. The challenge in friend recommendations lies in identifying users who might be interested in connecting with one another. While some connections are straightforward (i.e, people who went to the same school or work at the same company), many potential connections are hidden within the vast social graph. Graph polynomials offer a mathematical approach to uncovering these hidden connections. An adjacency matrix is a square matrix that represents a graph's connections. In the context of a social graph, the adjacency matrix will have a 1 in the cell (i, j) if there is a friendship between user i and user j , and 0 otherwise. The adjacency matrix polynomial is formed by taking powers of this matrix. The adjacency matrix polynomial is a tool used by Facebook to understand the structure of the social graph better. By raising the adjacency matrix to various powers and analyzing the resulting matrix, Facebook can extract valuable information about user relationships. For instance, they can identify transitive relationships, (i.e, "friends of friend"), even if these connections

are not explicitly represented in the graph. One of the key insights gained through the use graph polynomials is the concept of triadic closure. Triadic closure is a social theory that suggests that if user A is friends with user B and user A is friends with user C, then there's a higher likelihood that user B and user C will become friends, even if they don't know each other. This is a common phenomenon in social networks and can be leveraged for friend recommendations. Facebook uses graph polynomials to identify patterns of triadic closure in the social graph. They can calculate how closely connected two users are indirectly through a series of mutual friends. By assigning scores based on the strength of these indirect connections, Facebook can make more accurate friend recommendations. For example, if you are friends with person A, and person A is friends with person B and person C, Facebook's algorithms can identify the potential connection between you and person B and person C. The strength of this potential connection might be based on factors like the number of mutual friends, the frequency of interactions, and shared interests. Facebook's recommendation systems are powered by machine learning algorithms, which incorporate graph polynomials as one of the many features used for predicting potential friendships. Machine learning models take into account various data points, not just graph polynomials, to provide recommendations. These data points include user activity, location, interests and more. It's important to note that Facebook takes user privacy seriously. While they use complex algorithms to make friend recommendations, they also have stringent privacy policies and controls in place. Users can control who can see their information and who can send them friend requests. The use of graph polynomials for recommendations does not compromise user privacy. Friend recommendations are a complex problem, and there's always room for improvement. Facebook continually refines its algorithms to provide better suggestions. Challenges include dealing with fake profiles, improving recommendation diversity, and addressing issues related to algorithmic bias. The use of graph polynomials for friend recommendations also raises ethical corners related to user manipulation and filter bubbles. There's a fine balance between providing users with relevant recommendations and avoiding the creating of echo chambers or reinforcing existing biases. Facebook, like other social

media platforms, is under constant scrutiny in this regard. Facebook's use of graph polynomials for friend recommendations is a fascinating application of mathematics and machine learning. By leveraging the power of graph theory, they can identify hidden connections within the vast social graph and provide users with more relevant and valuable friend recommendations. These recommendations, powered by machine learning, play a crucial role in enhancing the overall user experience on the platform. However, it's important to remember that while these algorithms are powerful, user privacy and ethical considerations remain paramount.

4.2 Case Study 2: Drug discovery through graph polynomial analysis of molecular graphs

The process of drug discovery is a complex and multifaceted endeavor that involves the identification, design, and development of new therapeutic agents to treat various diseases and conditions. Over the years, advances in computational chemistry and bioinformatics have played a crucial role in accelerating the drug discovery process. One such innovative approach is the use of graph polynomial analysis of molecular graphs. Molecular graphs are mathematical representations of chemical compounds, where atoms are depicted as nodes and chemical bonds as edges. These graphs provide a structural foundation for understanding the relationships between atoms within a molecule. In drug discovery, molecular graphs are indispensable, as they allow researchers to visualize and analyze the structural properties of compounds, making it easier to identify potential drug candidates. By applying graph theory concepts, researchers can extract valuable information from molecular structures. This information includes connectivity, molecular symmetry, aromaticity, and more. A prominent aspect of graph theory in drug discovery is the application of graph polynomials, which are used to compute various molecular descriptors. The most commonly used graph polynomials in drug discovery are the characteristic polynomials, the chromatic polynomial, and the Hosoya polynomial. The characteristic polynomial of a molecular graph is used to calculate eigenvalues, which offer information about

molecular properties, such as molecular weight, molecular connectivity, and aromaticity. By analyzing the eigenvalues, researchers can compare different molecules and make informed decisions about drug candidates. The chromatic polynomial is primarily used to determine the number of ways a molecule can be colored with a certain number of colors without adjacent atoms having the same color. This concept is particularly useful in the study of molecular isomerism and the design of drugs with specific properties. The Hosoya polynomial focuses on the enumeration of subgraphs within a molecular graph, which is important in understanding the topological structure of molecules. It helps in identifying structural motifs that are essential for the activity of a compound. Molecular descriptors are numerical values derived from the analysis of molecular graphs using graph polynomials. These descriptors play a vital role in drug design and can be categorized into two main types: topological and physicochemical descriptors. Topological descriptors are derived from the structural properties of a molecule. Graph polynomials, such as the Hosoya polynomial and the chromatic polynomial, provide information for the calculation of topological descriptors. These descriptors can be used to predict a molecule's bioactivity and its interaction with biological targets. Physicochemical descriptors are derived from the physical and chemical properties of a molecule. They include properties such as molecular weight, logP (partition coefficient), polar surface area, and more. These descriptors are essential for understanding a molecule's solubility, pharmacokinetics, and potential toxicities. Graph polynomial analysis of molecular graphs is an integral part of virtual screening and molecular modeling. Virtual screening involves the rapid evaluation of a large database of chemical compounds to identify potential drug candidates. By using molecular descriptors derived from graph polynomial analysis, researchers can filter and prioritize compounds for experimental testing. This significantly reduces the time and resources required for drug discovery. Molecular modeling, on the other hand, encompasses a wide range of techniques for simulating and predicting the behavior of molecules. It relies on accurate representations of molecular structures, and graph polynomial analysis contributes to the creation of these representations. Molecular dynamics simulations, docking studies, and quanti-

tative structure-activity relationship (QSAR) modeling all benefit from the structural insights gained through graph polynomial analysis. Structure-Activity Relationship (SAR) studies are crucial in drug discovery, as they aim to establish a connection between the chemical structure of a compound and its biological activity. Graph polynomial analysis allows researchers to extract valuable topological information from molecular graphs, which be used to identify essential structural features responsible for a compound's activity. By comparing the molecular graphs and associated descriptors of active and inactive compounds, SAR studies help researchers optimize and modify molecules to enhance their activity, selectivity, and safety profiles. This iterative process is central to the development of effective drugs. Fragment-based drug design is a strategy that involves the systematic exploration of chemical fragments to identify potential drug candidates. Molecular graphs and graph polynomial analysis are instrumental in this approach. Researchers can break down complex molecules into smaller fragments and evaluate their structural and topological properties. This process aids in the identification of essential pharmacophores and the assembly of fragments into lead compounds. Graph polynomial analysis also plays a role in target identification, an essential step in drug discovery. By analyzing the interaction of molecular graphs and protein structures, researchers can identify potential drug targets. Graph-based approaches, such as ligand-protein interaction networks, are used to predict protein-ligand binding affinities and understand the underlying molecular interactions. While graph polynomial analysis of molecular graphs offers numerous advantages in drug discovery, it is not without challenges and limitations. Some of the key limitations include: Computational Intensity: Graph polynomial calculations can be computationally intensive, especially for large molecular graphs. This may limit the scalability of the approach. Data Availability: The quality and availability of data for training machine learning models based on graph polynomial descriptors can impact the accuracy of predictions. Interpretability: Interpreting the results of graph polynomial analysis and relating them to specific biological activities can be challenging. Hybrid Approaches: Combining graph polynomial analysis with other computational methods, such as molecular dynamics stimulations, and quantum

chemistry, is often necessary for a more comprehensive understanding of molecular behavior. Graph polynomial analysis of molecular graphs is a powerful tool in drug discovery. It provides a structured and quantitative approach to understanding the structural and topological properties of molecules. This methodology aids in virtual screening, molecular modeling, SAR studies, and fragment-based drug design, accelerating the drug discovery process. While it has limitations, ongoing research and innovation in this field are likely to address these challenges and further improve the accuracy and efficiency of drug discovery. The future of drug discovery through graph polynomial analysis holds the promise of more effective and personalized medicines for a wide range of diseases and conditions.

4.3 Case Study 3: Traffic flow optimization in urban planning

Traffic flow optimization in urban planning is a complex and critical task that involves managing and improving the efficiency of transportation networks within cities. Graph polynomials, specifically traffic flow models based on graph theory, play a crucial role in understanding and optimizing these systems. Traffic congestion is a major concern in urban areas, leading to wasted time, increased pollution, and decreased quality of life. To address these issues, urban planners and transportation engineers employ a variety of strategies, with one of the key components being traffic flow optimization. This optimization seeks to improve the movement of vehicles, pedestrians, and other forms of transportation in cities. In the context of traffic flow optimization, road networks can be effectively modeled as graphs. In this representation, each intersections or node corresponds to a point in the graph, and the roads connecting them form the edges. Graphs provide a powerful framework for analyzing and optimizing urban transportation networks. In traffic flow optimization, several graph polynomials are used to model and analyze transportation networks. Adjacency Matrix: An adjacency matrix is a fundamental graph polynomial that represents the connections between nodes in a graph. In the context of urban planning, this matrix

can be used to understand the connectivity of roads in a city and how traffic can flow between different areas. Incidence Matrix: The incidence matrix helps determine how many roads meet at a given intersections, which is vital for optimizing traffic signals and managing congestion. Laplacian Matrix: The Laplacian matrix describes the behavior of a graph in terms of its eigenvalues and eigenvectors. It is useful for understanding the network's stability and efficiency in traffic flow. Flow Polynomials: Flow polynomials are used to represent traffic flow patterns in a network. They help model the movement of vehicles through road networks and are essential for optimizing traffic signal timings and route planning. Urban planners face numerous challenges in optimizing traffic flow. Some of the most pressing issues include: 1. Congestion : Congestion is a major problem in cities, resulting in traffic jams and increased travel times. Optimizing traffic flow aims to reduce congestion and improve the overall efficiency of the transportation network. 2. Environmental Impact: Traffic congestion contributes to increased pollution and carbon emissions. Reducing congestion through optimization can lead to environmental benefits. 3. Safety: Safety is a significant concern in urban traffic flow. Efficient optimization should prioritize the safety of all road users, including pedestrians and cyclists. 4. Data Collection: Gathering accurate and up-to-date data is essential for effective traffic flow optimization. This data includes traffic counts, road conditions, and real-time information. Graph polynomials provide a mathematical foundation for addressing these challenges in traffic flow optimization: 1. Optimal Signal Timing: Graph polynomials can be used to model traffic signal timings. By optimizing these timings, urban planners can improve the overall flow of traffic and reduce congestion. 2. Route Planning: Graph algorithms, such as Dijkstra's algorithm or A* search, can be employed to find the most efficient routes for vehicles, minimizing travel times and reducing congestion. 3. Network Design: Graph theory helps in designing transportation networks that are efficient and minimize bottlenecks. By analyzing the structure of the graph, planners can make informed decisions about road expansions or new infrastructure. 4. Traffic Simulation: Traffic simulation models based on graph polynomials allow planners to predict and evaluate the impact of changes to the transportation network, such as the

addition of new roads or public transit systems. 5. Public Transportation Optimization: Optimizing the flow of public transportation, including buses and subways, is critical in reducing traffic congestion. Graph polynomials can assist in designing efficient public transportation routes. Several cities around the world have successfully used graph theory and graph polynomials in their traffic flow optimization efforts;

1. Los Angeles, USA; Los Angeles employed traffic flow optimization strategies that included optimizing traffic signal timings, implementing intelligent transportation systems, and using graph based models to improve freeway capacity.
2. Singapore: Singapore uses dynamic traffic management systems that rely on graph polynomials to optimize traffic flow. The city-state has also integrated public transportation networks with road networks for seamless urban mobility.
3. Amsterdam, Netherlands: Amsterdam is known for its efficient transportation system, and graph-based models have played a significant role in optimizing traffic flow, reducing congestion, and promoting cycling and public transportation. Traffic flow optimization in urban planning is essential for creating sustainable, efficient, and safe transportation networks. Graph polynomials rooted in graph theory, provide a valuable framework for modeling, analyzing, and optimizing these complex systems. By using mathematical models and real-time data, urban planners can make informed decisions to address congestion, reduce environmental impact, enhance safety, and improve the overall quality of life in cities. As technology and data analysis techniques continue to advance, the role of graph polynomials in traffic flow optimization will become increasingly critical in shaping the future in urban mobility.

Chapter 5 Conclusion and Future Research

The future of graph polynomials is bright and promising. Graph polynomials represent a fascinating area of study at the intersection of graph theory, algebra and combinatorics. These polynomials have found applications in wide range of fields. As we look to the future, several exciting developments and expectations can be highlighted regard the role of graph polynomials in various scientific and practical domains. Graph polynomials have played a significant role in network analysis, where they provide insights into the structural properties of networks. In the future, we can expect an even wider range of applications, such as in social networks, transportation networks, and biological networks. As our world becomes increasingly interconnected, understanding and modeling these complex systems will be of importance. As the importance of secure and resilient networks continues to grow, we can anticipate further developments in using graph polynomials for intrusion detection and threat analysis. With the emergence of quantum computing as a transformative technology, graph polynomials are expected to play a crucial role in algorithm development and quantum network analysis. When it comes to machine learning and data science, graph polynomials can be employed to represent data that make it suitable for various machine learning tasks. In the future, we can expect the integration of graph polynomials in more advanced machine learning algorithms, improving our ability to analyze and model complex datasets. In algebraic graph theory, graph polynomials will continue to be at the core of this field, helping researchers uncover deeper connections between algebra and graph theory. With ongoing advancements in computational power and algorithmic techniques, we can anticipate more efficient algorithms for calculating graph polynomials. This will make it easier to analyze larger and more complex graphs, further expanding their applicability. Researchers

from various fields will work together to harness the power of graph polynomials to address complex problems, leading to groundbreaking discoveries and applications. As graph polynomials become increasingly important in various domains, there will be a growing need for education and outreach. Universities and research institutions are likely to offer specialized courses and resources to train the next generation of researchers and professionals in this field.

In conclusion, the future of graph polynomials is bright and promising. Their versatility and power in modeling and analyzing complex structures make them very valuable to a wide range of applications. As technology and research continue to advance, graph polynomials are expected to play an ever-expanding role in solving real-world problems, making them a critical area of study for mathematicians, scientists and engineers.

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